Modeling Heterogeneous Semantics in Ptolemy

with Modular Actor Interfaces

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Heterogeneous Systems



Different parts of a system and different levels of hierarchy can have different Models of Computation



Reasoning About Heterogeneity

- There needs to be a way to reason about the composition of models with different Models of Computation.
- This is necessary for establishing a consistent semantics for heterogeneous systems.
- This can be used to do analyses and forms of verification that cross the boundaries of heterogeneous composition.

Ptolemy II



- An open source research platform for modeling systems.
- Implements an Actor Oriented language for model development.
- Allows hierarchical design and the heterogeneous composition of different Models of Computation.
- ~2.5 Million Lines of Code.
- Code generation.

Actors



Some Specific Actors



Actor Graphs



Author: Gang Zhou (based on Ptolemy Classic demo)

Models



Author: Gang Zhou (based on Ptolemy Classic demo)

Dataflow Directors : DDF, PSDF, CSDF, MDSDF, HDF, SDF, etc...



SDF : Production and Consumption rates are fixed. Scheduability is decidable, and if a model is scheduable, it can be scheduled periodically with fixed-sized buffers.

SR, DE, CT, etc...



SR : If actors in a model are monotonic then there exists a unique solution for the value of every relation that can be reached through a fixed-point iteration.

Composite Actors



Modular Actor Interface











Updates to a new state from inputs and state

Deadline

Determines a deadline for execution from inputs and state





Timeupdate

Updates to a new state from inputs, state, and time



Directors compose the interfaces of Actors in a Model into those of the Composite Actor.

Mathematical Representations

The fixed-point semantics of SR can be understood in clear mathematical terms.

$$(\widetilde{F}_{s,x}(y))(o) = F_j(s_j,(x,y) \upharpoonright_{I_j})(o)$$

Letting $y_{s,x}^*$ be the unique least fixpoint of $\widetilde{F}_{s,x}$.

$$F(s,x) = y_{s,x}^* \upharpoonright_O P(s,x) = \left(P_1(s_1, (x, y_{s,x}^*) \upharpoonright_{I_1}), ..., P_n(s_n, (x, y_{s,x}^*) \upharpoonright_{I_n}) \right)$$

From "A Modular Formal Semantics for Ptolemy", Tripakis et al.

Coroutine Model of Computation

Theorem 1. Given a non-strict Coroutine Model \mathcal{M} , if the input $\mathbb{I}_{\mathcal{M}}$ and output $\mathbb{O}_{\mathcal{M}}$ types of the model are finite-height pCPOs and operator \oplus is monotonic, then the above recursive equations characterizing the kernel functions \mathbf{e} and \mathbf{f} have unique least fixed-point solutions in the partial order of functions with codomains $2^{\mathbb{G}}$ and $\mathbb{O}_{\mathcal{M}}$, respectively.

Theorem 2. Given a non-strict Coroutine Model \mathcal{M} , if the input $\mathbb{I}_{\mathcal{M}}$ and output $\mathbb{O}_{\mathcal{M}}$ types of the model are finite-height pCPOs and operator \oplus is monotonic, and if for each $q \in Q$ the functions enter_q and fire_q are monotonic in terms of \mathbb{I}_q , and the mapping functions $m_{\mathbb{I}}$ and $m_{\mathbb{O}}$ are monotonic, then the non-strict kernels \mathbf{e} and \mathbf{f} are continuous in terms of $\mathbb{I}_{\mathcal{M}}$.

"The Coroutine Model of Computation", Shaver and Lee, 2012

I Conclude !

- Try out Ptolemy! It's open and free!
- Check out Modular Actor Interfaces.
- Feel free to design your systems modularly, hierarchically, and heterogeneously.
- <u>http://ptolemy.eecs.berkeley.edu</u>